

Kant: Logic is a science of the necessary laws of thought.

Logic is the anatomy of thought. Locke

Logic is the art of going wrong with confidence. Joseph Blatch.

L1-1

1.1

L1: LOGIC

We'll need for 3.1.

Gateway to precise thinking

A proposition is a statement that is either true (T) or false (F).

Not tested

truth values
Ex

Are props.

Are not props.

All Gore is a Dem. (T)

G.W. Bush is a Dem. (F)

The sky is blue.

(but depends on
context, opinion)

Whazzup?

Just do it.

right, Texas

$1+1=2$, (T)

$1+1=3$, (F)

$2x=10$ if $x=5$. (T)

$2x=10$.

[assume we know
nothing about x]

In CS, 0 = false, 1 = true.

In fuzzy logic (AI), a proposition can have
a truth value between 0 and 1.

Ex Ken is tall. (0.3)

Overhead

Japan
Quantum meth
bit/fig (Kashi 24)

Let p, q, r , etc. denote propositions.
We use logical operators to build compound propositions.

① $\neg p$ means $\begin{cases} \text{not } p \\ \text{the negation of } p \\ \text{it is not the case that } p \end{cases}$

Truth table for $\neg p$

p	$\neg p$
T	F
F	T

If p is a $\begin{matrix} T \\ F \end{matrix}$ prop., then $\neg p$ must be a $\begin{matrix} F \\ T \end{matrix}$ prop.

Ex $p: 1+1=2$ (T)
 $\neg p: 1+1 \neq 2$ (F)

Ex a, b are real #s

$p: a > b$

$\neg p: \text{it is not the case that } a > b$
(i.e., $a \leq b$)

$\begin{cases} \swarrow \text{one is T;} \\ \searrow \text{the other} \\ \text{is F} \end{cases}$

Note: \neg is a unary operator.

Connectives are operators that combine
2 or more props.

② $p \wedge q$ means $\begin{cases} p \text{ and } q \\ \text{the conjunction of } p \text{ and } q \end{cases}$

Truth table

p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

all possib

$p \wedge q$ is a T prop. $\xleftrightarrow{\text{exactly when}}$
 p and q are both T.

Ex p : Gore will win. (?)
 q : $1+1=3$. (F)
 $p \wedge q$: Gore will win and $1+1=3$. (F, regardless)

Your grandma
 went w/ my
 bob.

③ $p \vee q$ means $\left\{ \begin{array}{l} p \text{ or } q \text{ (inclusive or)} \\ \text{the disjunction of } p \text{ and } q \end{array} \right.$

Truth table

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

} one or both of p, q T

$p \vee q$ is a F prop. \iff
 p and q are both F

④ $p \oplus q$ means $\left\{ \begin{array}{l} \text{exclusive or of } p \text{ and } q \\ p \text{ xor } q \end{array} \right.$

Truth table

p	q	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

\leftarrow different from $p \vee q$

for purposes of circuit design, it's interesting that $p \oplus q$ is F \iff p, q have same truth value

$p \oplus q$ is a T prop. \iff
 exactly one of p or q is T (but not both!)
 ("soup or salad" ideal)

Ex Let's say

p : You have soup. (T)

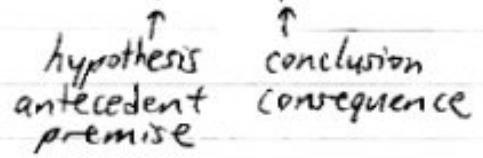
q : You have salad. (T)

Then, $p \oplus q$: You have soup xor salad. (F)

but $p \vee q$: (T)

Note: In math, "or" is an inclusive or (3), unless otherwise specified.

⑤ The implication $p \rightarrow q$



SNL sketch

Ex d : Pat is a Daddy.

m : Pat is a man.

There are many ways to say $d \rightarrow m$

- if d , [then] m
- d implies m
- d only if m
- d is sufficient for m

- m if d
- m whenever d
- m is necessary for d

$(p \rightarrow q \text{ is a F prop}) \iff (p \text{ is T and } q \text{ is F})$
 "Breaking a contract"

Truth table

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

} If the hypothesis (p) is F, the prop. $p \rightarrow q$ is T "by default".

Ex $p: 1+1=3$ (F)

$q: \text{Gore will win}$ (?)

Then, $p \rightarrow q: \text{If } \underbrace{1+1=3}_F, \text{ then } \underbrace{\text{Gore will win}}_?$ (T "by default")

Ex If $\underbrace{1+1=2}_T$, then $\underbrace{\text{a lemon is yellow}}_T$ (all we need) (T)

p and q need not be related! (stats: correlation does not imply causation. elev.)

Ex If $\underbrace{\text{Pat is a Daddy}}_p \text{ (or d)}$, then $\underbrace{\text{Pat is a man}}_q \text{ (or m)}$.

This prop. must be T, because (realistically) there's no way p is T and q is F.

Math Ex If an integer ends in a "3", it's an odd #.

Same idea!

Key: q is T
 Math produce

Why does this sound weird?
 no causal rel-ship

Assume Pat is human.
 proper name

Math

4.T, Fermat
 $x^{2m} + y^{2m} = z^{2m}$, p11.
use no prime power
integer sol)

Think of an "it-then" statement as a conjecture (unproven hunch)
If the prop is T, we have a theorem we can put in books.
If the prop. is F, no grant \$\$\$ sink!

5 More

The converse of $p \rightarrow q$ is $q \rightarrow p$.
The contrapositive of $p \rightarrow q$ is $\neg q \rightarrow \neg p$.

not equivalent
These are "logically equivalent" (Sec 1.2):
both are T or both are F.

HW 1.2 #16

depending on
whether the
truth values
of p, q

Ex Write the converse of
"Pat is a man whenever Pat is a Daddy."

Rewrite in "it-then" form:
"If Pat is a Daddy, then Pat is a man."
P q

The converse is:
"If Pat is a man, then Pat is a Daddy."
q P

This new prop. is fundamentally different
from the old one!

Ex The contrapositive is:
"If $\neg q$, then $\neg p$ "
"If Pat is not a man, then
Pat is not a Daddy."

This is fundamentally the same as
the old prop.!

⑥ The biconditional $p \leftrightarrow q$

means $\left\{ \begin{array}{l} p \text{ if and only if } q \\ p \text{ is necessary and sufficient for } q \\ \text{if } p \text{ then } q, \text{ and if } q \text{ then } p \\ \text{"conversely"} \end{array} \right.$

$(p \leftrightarrow q \text{ is } T) \iff (p, q \text{ are both } T \text{ or both } F)$

Truth table

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

We tend to look at the table from right to left.

Ex If we know this is a true prop:
"Pat is a man \leftrightarrow Pat has an Adam's apple," then
both are T, or both are F

HW Tips (Relevant: pp. 1-8; ignore Boolean searches, bit strings)

Ex p : I am young.
 q : I am stupid.
 r : I can attend the movie.

Express $(p \vee q) \rightarrow \neg r$ as an English sentence.

Order of ops: $()$ 1st, \neg next, \vee/\wedge , $\rightarrow/\leftrightarrow$

"If I am young or stupid, then I cannot attend the movie."

Ex (23d) Construct a truth table for $(p \rightarrow q) \wedge (\neg p \rightarrow q)$

p	q	$(p \rightarrow q)$	$\neg p$	q	$(\neg p \rightarrow q)$	\checkmark
T	T	T	F	T	T	T
T	F	F	F	F	T	F
F	T	T	T	T	T	T
F	F	T	T	F	F	F

Different approaches ok!

Concise answer:

p	q	$(p \rightarrow q) \wedge (\neg p \rightarrow q)$
T	T	T
T	F	F
F	T	T
F	F	F

(9a) Converse:

If I will ski tomorrow, then it snows today.

Let's say both are T.

Then, this implication is T.

In conversation, a cause-and-effect relationship is implied.

Our sense of logic does not focus on content or cause-and-effect.