

### 3.1: METHODS OF PROOF

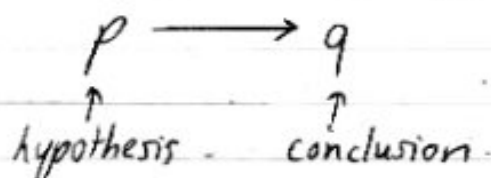
A conjecture is a proposition whose truth value (T/F) is unknown. If it can be proven to be T, it becomes a theorem.

A lemma is a "smaller pre-theorem" used to prove a "larger" theorem.

!kōr-ə, lər-ē  
!hāv  
Brit: kə-'räl-ərə

A corollary is a "post-theorem" that follows directly from a previous theorem.

Many conjectures can be written as implications



How can I  
rewrite this  
as an if-then  
stmt?

Ex "All primes are odd" can be rewritten as  
 "If  $\underbrace{n \text{ is prime}}_p$ , then  $\underbrace{n \text{ is odd}}_q$ ."

When  $n=2$ ,  $p$  is T and  $q$  is F.  
 So,  $n=2$  is a counterexample, and  
 the conjecture is F.

Technically:  $\forall x \underbrace{O(x)}_{\substack{\text{(primes)} \\ x \text{ is odd}}}$  is F.

If such a conjecture is T, we have the theorem

$$p \implies q$$

A proof of a theorem can include:

① Axioms / Postulates

- statements accepted as true

Ex  $0 \neq 1$

underlying  
assumps  
about math.  
structures

② Other Theorems / Lemmas

③ Definitions

Ex  $n$  is an even number  $\Leftrightarrow$   
 $\exists k \in \mathbb{Z} : n = 2k$

Ex  $n$  is an odd number  $\Leftrightarrow$   
 $\exists k \in \mathbb{Z} : n = 2k + 1$

Another  
characins  
mod 2 arithmetic  
even  $\Leftrightarrow 2|n$

④ The hypotheses of the theorem

- assume  $p$  is true (show  $q$  must be true)  
- used in direct proofs, proofs by cases

⑤ The negation of the conclusion

- assume  $\neg q$  is true (show  $\neg p$  must be true)  
- prove  $p \rightarrow q$  is  $\vdash$  by proving that the  
contrapositive  $\neg q \rightarrow \neg p$  is  $\vdash$ .  
- used in indirect proofs

## ⑥ Rules of Inference

- logical rules (Table 1, p.169)  
(Table 2, p.174)
- basic (don't need to memorize names)

Ex A lemma states " $r \Rightarrow s$ " ( $r \rightarrow s$  is T)  
We know  $r$  is T.  
Therefore,  $s$  is T.

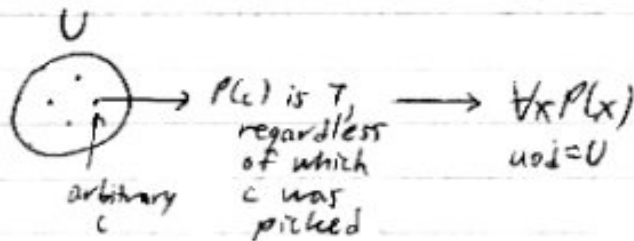
Shorthand:

$$\begin{array}{c} r \\ \hline r \rightarrow s \\ \hline \therefore s \\ \uparrow \\ \text{therefore} \end{array}$$

"modus ponens" rule  
critical in induction  
(i.e.,  $[r \wedge (r \rightarrow s)] \rightarrow s$  is a tautology)

## Ex "Universal Generalization"

$P(c)$  for an arbitrary  $c \in U$   
 $\therefore \forall x P(x)$



## FALLACIES

- incorrect inferences

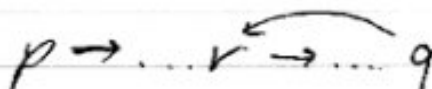
Affirming the  
Consequent  
John McLaughlin  
 $p \rightarrow q$  right

Ex  $(p \rightarrow q) \Leftrightarrow (q \rightarrow p)$  Usually wrong!  
converse

Ex  $(p \rightarrow q) \Leftrightarrow (\neg p \rightarrow \neg q)$  Usually wrong!  
inverse

Descartes'  
proof of  
 $\exists$  God

Ex Begging the question / Circular reasoning  
- when a theorem is used to prove itself  
- for example, when proving  $p \Rightarrow q$  directly, a step in the proof relies on  $q$  being true.



TYPES OF PROOFS

## ① Direct Proofs

Ex (#22) Prove that the product of two rational #s is rational. ( $\mathbb{Q}$  = set of all rational #s)

(Optional Rewrite)

If  $m$  and  $n$  are rational #s,  
then  $mn$  is rational.

Assume / Suppose  $m$  and  $n$  are rational #s. ( $m, n \in \mathbb{Q}$ )

(By definition)

$$m = \frac{a}{b} \quad \text{and} \quad n = \frac{c}{d}$$

for some  $a, b, c, d \in \mathbb{Z}$  where  $b, d \neq 0$ .

(Sometimes, we assume that the fractions are in reduced form, but not here.)

$$\text{Then, } mn = \left(\frac{a}{b}\right)\left(\frac{c}{d}\right) = \frac{ac}{bd}$$

where  $ac \in \mathbb{Z}$ ,  $bd \in \mathbb{Z}$ , and  $bd \neq 0$ .

$\therefore mn$  is rational ( $mn \in \mathbb{Q}$ )

QED

Ex 14 (p. 175)  $n$  odd  $\Rightarrow n^2$  odd

quod erat  
demonstrandum

Test: odd\*even = odd

See (3-3.5) ② Indirect Proofs

Ex 15, p. 175  $3n+2$  odd  $\Rightarrow n$  odd

Ex If  $\underbrace{x \in \mathbb{Z}}_{p_1}$  and  $\underbrace{y \in \mathbb{Z}}_{p_2}$ , then  $\underbrace{2x+2y \neq 15}_q$   
 $p = p_1 \wedge p_2$

Prove the contrapositive:  $\neg q \rightarrow \neg p$

Assume  $\neg q$ :  $2x+2y=15$

$$\Rightarrow x+y = \frac{15}{2}$$

$$\Rightarrow \underbrace{x \notin \mathbb{Z} \text{ or } y \notin \mathbb{Z}}_{\neg p}$$

Too close  
to  $\neg p$

Note: If  $p = p_1 \wedge p_2 \wedge \dots$  (conjunctive hypothesis)  
to show  $\neg p$  is T,  
it suffices to show that one of  
 $p_1, p_2, \dots$  must be F

## ② Indirect Proofs

Direct  
mod 2 proof:  
 $7n-11 \equiv 1 \pmod{2}$   
 $7n \equiv 12 \pmod{2}$   
 $n \equiv 0 \pmod{2}$

Ex Prove: If  $7n-11$  is odd, then  $n$  is even.

$\underbrace{7n-11 \text{ is odd}}_p, \quad \underbrace{\text{then } n \text{ is even}}_q$

Indirect  
mod 2 proof:  
 $n \equiv 1 \pmod{2}$   
 $7n \equiv 7 \pmod{2}$   
 $7n-11 \equiv 1 \pmod{2}$   
 $7n-11 \equiv -10 \pmod{2}$   
 $7n-11 \equiv 0 \pmod{2}$   
 $2 \mid 7n-11$

Prove the contrapositive:  $\neg q \rightarrow \neg p$   
i.e.,  $n$  is odd  $\rightarrow 7n-11$  is even

$\uparrow$   
( $n$  alone is easier to start with)

Assume  $\neg q$ :  $n$  is odd.

Then,  $\exists k \in \mathbb{Z}: n = 2k+1$

$$\begin{aligned} \Rightarrow \exists k \in \mathbb{Z}: 7n-11 &= 7(2k+1)-11 \\ &= 14k+7-11 \\ &= 14k-4 \\ &= 2(7k-2) \end{aligned}$$

$\in \mathbb{Z}$

$\Rightarrow 7n-11$  is even ( $\neg p$ )

QED

### ③ Proofs by Contradiction / Reductio ad absurdum

Gauß's zero  
↓

To prove that a conjecture  $c$  is true,  
show that  $\neg c$  leads to a contradiction ( $*$ ).

i.e., Show that, for some proposition  $r$ ,

$$\neg c \rightarrow (r \wedge \neg r) \text{ is } T$$

$$\underbrace{\hspace{10em}}_{\text{must be } F}$$

So,  $\neg c$  must be  $F$ .  
 $c$  must be  $T$ .

Ex 18, pp. 176-7  $\sqrt{2}$  is irrational ( $\sqrt{2} \notin \mathbb{Q}$ )

Ex If  $\underbrace{x \in \mathbb{Z} \text{ and } y \in \mathbb{Z}}_p$ , then  $\underbrace{2x + 2y \neq 15}_q$ .

Assume  $p: x \in \mathbb{Z}$  and  $y \in \mathbb{Z}$

Assume  $\neg q: 2x + 2y = 15$

$$\Rightarrow x + y = \frac{15}{2}$$

$$\Rightarrow x + y \notin \mathbb{Z}$$

$\left. \begin{array}{l} \Rightarrow x + y \notin \mathbb{Z} \\ \Rightarrow x + y \notin \mathbb{Z} \end{array} \right\} r \text{ and } \neg r *$

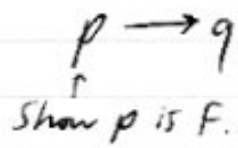
However,  $p \Rightarrow x + y \in \mathbb{Z}$  (sum of two integers  $\in \mathbb{Z}$ )

So, assuming  $p$  is  $T \Rightarrow \neg q$  is  $F \Rightarrow q$  is  $T$



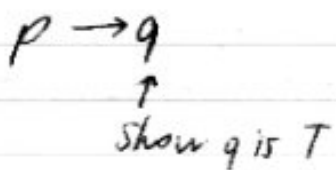
④ Vacuous proof

Show  $P(n)$  true  
if  $P(n); (n \geq 1) \rightarrow q$



don't  
worry

⑤ Trivial proof



Proof by Cases

Ex Prove  $\forall n \in \mathbb{Z} P(n)$

Can prove  $\forall n \in \mathbb{Z}^{\text{even}} P(n)$  "Case 1"  
and  $\forall n \in \mathbb{Z}^{\text{odd}} P(n)$  "Case 2"

Ex +, 0, -

Ex Congruence classes (mod m)

Ex  $x > y, x = y, x < y$

To prove  $p \leftrightarrow q$

Prove  $p \rightarrow q$  > not necessarily  
and  $q \rightarrow p$  directly

## EXISTENCE PROOFS

Ex  $\exists n P(n)$

Constructive existence proofs indicate how to find a value(s) that make(s)  $P$  true.

Sufficient to  
prove  $\forall n \in \mathbb{Z}^+$   
...  $n$ ...

Ex Prove that, for every  $n \in \mathbb{Z}^+$  ( $n \geq 2$ ), there is a sequence of  $n-1$  consecutive composite integers.

(This means there are arbitrarily large gaps between the primes.)

Consider: (Assume  $n$  big for now.)

$$\begin{array}{l} n! \leftarrow \begin{array}{l} \text{prime if } n=2 \\ \text{not if } n \geq 2 \end{array} \\ n!+1 \leftarrow ? \\ \text{sequence of } n-1 \text{ composites } \left\{ \begin{array}{l} n!+2 \leftarrow 2|n!, 2|2, \text{ so } 2|(n!+2) \\ n!+3 \leftarrow 3|n!, 3|3, \text{ so } 3|(n!+3) \\ \vdots \\ n!+n \leftarrow n|n!, n|n, \text{ so } n|(n!+n) \end{array} \right. \end{array}$$

$$\{(n!+k) \mid k=2, 3, \dots, n\}$$

Ex 24 (p. 180) gives a nonconstructive proof showing there are  $\infty$  many primes.

Proof by Contradiction

Assume there are only finitely many primes,  
say  $n$  of them.

Primes:  $p_1, p_2, \dots, p_n$  ← <sup>complete</sup> list of all primes  
 $\quad \quad \quad \parallel \quad \parallel$   
 $\quad \quad \quad 2 \quad 3$

Consider  $p_1 p_2 \dots p_n + 1 \leftarrow M$

$p_i \nmid M \quad (1 \leq i \leq n)$   
 because there is always  
 a remainder of 1 if  
 $M$  is divided by any  $p_i$   
 $(M \equiv 1 \pmod{p_i})$

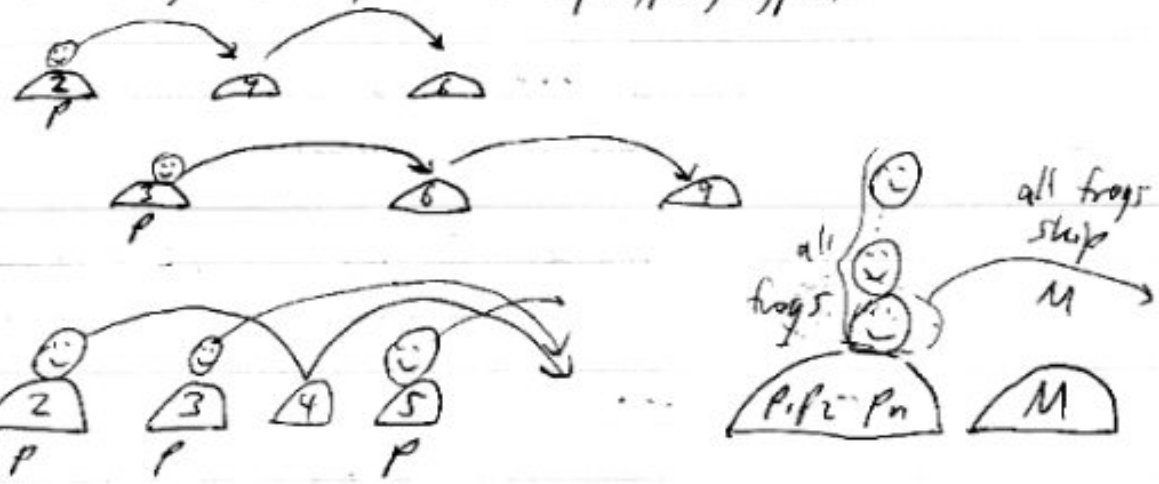
So, ①  $M$  itself is a new prime  
 not in the list. \*

or ②  $M$  is divisible by some prime  $(M$   
 that I missed. \*)

In either case, the completeness of my list is  
 contradicted.  
 QED

$M$  div'ed  
 by some  
 prime not  
 in my  
 list (maybe  
 $M$  itself)

Picture: Frogs correspond to  $p_1, p_2, \dots, p_n$



$p_1 \cdot p_n$  is  
lowest int. +  
divisible by  
all  $p_i$ ,  
by CRT,  
 $x \equiv 0 (p_1)$   
 $x \equiv 0 (p_n)$   
 $\vdots$   
 $\text{forall } x \equiv 0 (p_i p_n)$

Either  $M$  is prime



OR

A frog we missed lands on  $M$

