

6.1

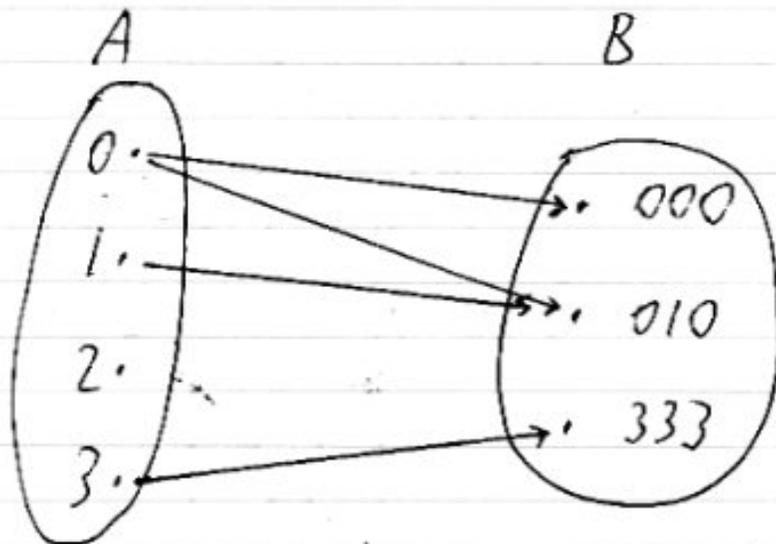
A binary relation (usually denoted by "R") from a set A to a set B is a subset of $A \times B$. It's a set of ordered pairs.

Ex $A = \{0, 1, 2, 3\}$

$B = \{000, 010, 333\}$

Let's say R is a relation from A to B s.t.
 $(a, b) \in R \iff$ digit a is in string b.

Picture of R:



Directed graph (digraph) - arrows or edges.

Some relations (not this one) are functions.
 Our R isn't a function

Why not?

0. ↗

2. (no arrow)

function



$$R = \{ (0,000), (0,010), (1,010), (3,333) \} \subseteq A \times B$$

Notation: $a R b \leftrightarrow (a,b) \in R$
 $a \not R b \leftrightarrow (a,b) \notin R$

Ex $1 R 010$

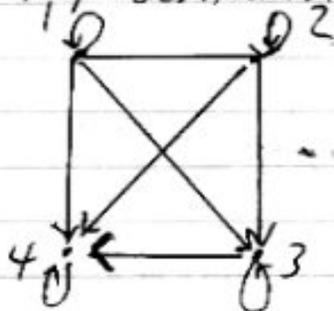
Special Case

R can be a relation from set A to itself.
 "R is a relation on A."

Ex $A = \{1, 2, 3, 4\}$

$R = \{ (a_i, a_j) \mid a_i \leq a_j \}$ is a relation on A

Digraph representation of R :



$$R = \{ (1,1), (1,2), (2,2), (1,3), (2,3), (3,3), (1,4), (2,4), (3,4), (4,4) \}$$

6.2: n-ary Relations

An n-ary relation on the sets: A_1, A_2, \dots, A_n is a subset of $A_1 \times A_2 \times \dots \times A_n$ ($n = \text{"degree"}$).

Ex (Database of students)

	Fields (A_i sets)			
	Last Name	First Name	ID #	Major
Records - elements of R	Jones	Paul	312	CS
	Smith	John	463	CS
	Smith	Paul	110	Math

The ID # field is a primary key, since there are no repeats. An ID # determines a 4-tuple.

6.3: Representing Binary Relations

- ① Digraphs
- ② 0-1 Matrices

old Ex

		B		
		00^10	01^20	3^333
A	$0 a_1$	1	1	0
	$1 a_2$	0	1	0
	$2 a_3$	0	0	0
	$3 a_4$	0	0	1

m_{ij} = entry in:
row i , column j ;
= $\begin{cases} 1 \leftrightarrow a_i R b_j \\ 0 \leftrightarrow a_i \not R b_j \end{cases}$

Old Ex

$$A = \{1, 2, 3, 4\}$$

$R = \{(a_i, a_j) \mid a_i \leq a_j\}$ is a relation on A

① Digraph ✓

② 0-1 Matrix Rep.

must be in same order

		1	2	3	4
1	1	1	1	1	1
2	0	1	1	1	1
3	0	0	1	1	1
4	0	0	0	1	1

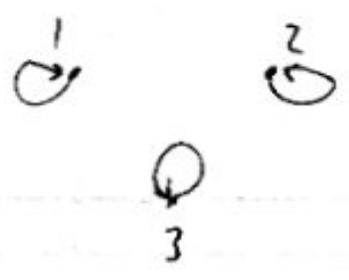
6.1/6.5: Equivalence Relations

Let's say R is a relation on a set A .

R is reflexive $\iff \forall a \in A (a, a) \in R$
"aRa"

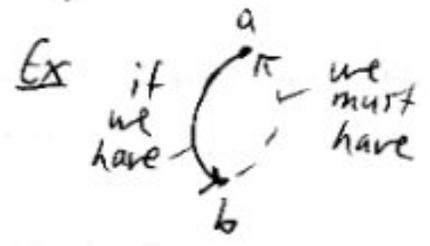
Ex $A = \{1, 2, 3\}$

R is reflexive \iff the corresponding digraph has at least more edges are OK



Property is implied in English usage.

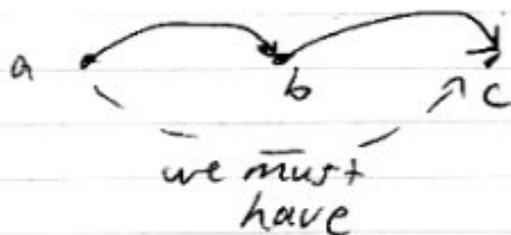
R is symmetric $\iff [\forall a, b \in A (aRb \rightarrow bRa)]$



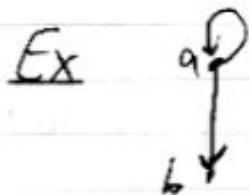
R is transitive \leftrightarrow

$$[\forall a, b, c \in A (aRb \text{ and } bRc \rightarrow aRc)]$$

Ex If we have



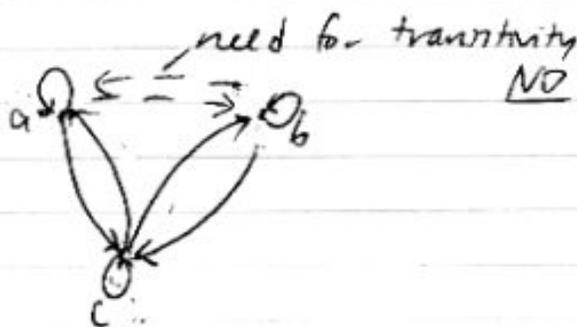
R is an equivalence relation \leftrightarrow
it is reflexive, symmetric, and transitive



NO
not reflexive (no $\begin{smallmatrix} a \\ a \end{smallmatrix}$)

not symmetric (we have $\begin{smallmatrix} a \\ b \end{smallmatrix}$, so we need $\begin{smallmatrix} b \\ a \end{smallmatrix}$)

Ex



In an EC, every elt is related to every other.

Ex

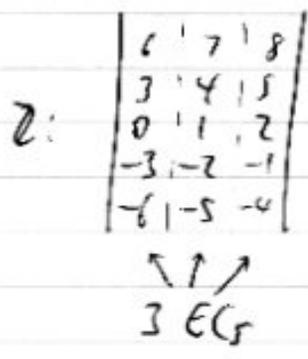


YES

(E.R.)

In an equivalence relation, the elements of an equivalence class are all related to one other and themselves.

Ex $A = \mathbb{Z}$
 $R = \{(a, b) \mid a \equiv b \pmod{3}\}$ is an E.R. on A .



Ex $A =$ freshmen at a college (forced into apts.)
 $R = \{(a, b) \mid a \text{ lives in the same apt. as } b\}$ is an E.R. on A .

ECs correspond to the apts.

Why is it R? S? T?

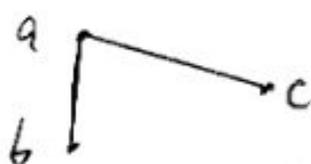
evil ban

CH 7: GRAPHS

A graph consists of dots (vertices) and edges that connect pairs of vertices.

$$\begin{aligned} \text{Ex } V &= \{a, b, c\} && \leftarrow |V| = 3 \text{ vertices} \\ E &= \{\{a, b\}, \{a, c\}\} && \leftarrow |E| = 2 \text{ edges} \end{aligned}$$

The graph $G = (V, E)$ can be represented as



a and b are adjacent vertices, because an edge connects them.
The edge $\{a, b\}$ is incident with a and b.

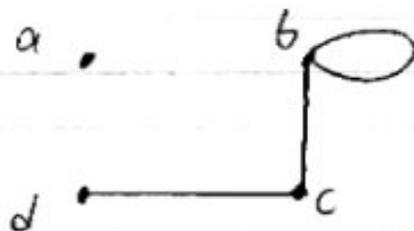
A simple graph has no loops or multiple edges.



The degree of a vertex " v " = " $\text{deg}(v)$ "

= # of edges incident with v
(a loop adds two to the degree).

Ex



$$\begin{aligned} \deg(a) &= 0, \text{ a is "isolated"} \\ \deg(b) &= 3 \\ \deg(c) &= 2 \\ \deg(d) &= 1 \end{aligned}$$

$$\text{sum of degrees} = 6$$

$$|E| = 3$$

Handshaking Thm

$$\sum_{v \in V} \deg(v) = 2|E|$$

degree sum = twice the # of edges

Why? Each edge contributes 2 to the degree sum.

$$\begin{array}{c} \text{---} \\ +1 \quad +1 \end{array} \quad 0+2$$

Corollary 1 The degree sum must be even.

Ex In a group of 5 computers, is it possible for

- 1 computer to be directly connected to 2 others and Another
- 4 others and
- 3 others and
- 3 others and
- 1 other.

Vertices = computers
 Edges = direct connections

Degree sum = 13 ← odd, so impossible!

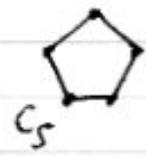
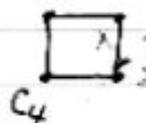
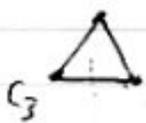
Corollary 2 A graph cannot have an odd number of odd-degree vertices.

Otherwise, the degree sum would be odd.

Special Graphs

Cycle graphs

C_n ($n \geq 3$)



Note: C_4 could look like

$|V| = n$
 $|E| = n$
 Each $\deg(v) = 2$

Complete graphs - every vertex pair is connected by exactly one edge
 (everyone shakes everyone else's hand once)

K_n



K_3



K_4



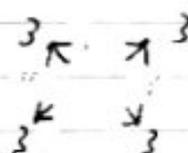
K_5

$|V| = n$

Each $\deg(v) = n - 1$

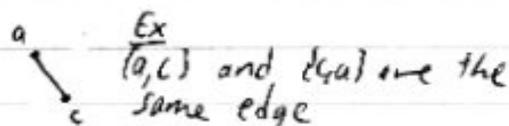
$|E| = ?$

Each of the n vertices is adjacent to $n - 1$ others.



Think: Incl-Excl
 Sort-of →
 familiar?

$$\frac{n(n-1)}{2}$$



How can we see that $\frac{n(n-1)}{2} = 1 + 2 + \dots + (n-1)$?

Ex $n=4$



3 edges



2 edges



1 edge

$|E \text{ for } K_4| = 3 + 2 + 1$

$|E \text{ for } K_n| = (n-1) + (n-2) + \dots + 1$

Proofs for: $1+2+\dots+n = \frac{n(n+1)}{2}$

$$\begin{aligned}
 1) \quad S &= 1+2+\dots+n \\
 S &= n+(n-1)+\dots+1 \\
 2S &= (n+1)+(n+1)+\dots+(n+1) \\
 2S &= n(n+1) \\
 S &= \frac{n(n+1)}{2}
 \end{aligned}$$

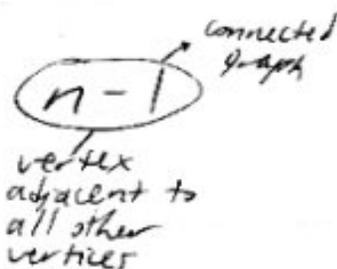
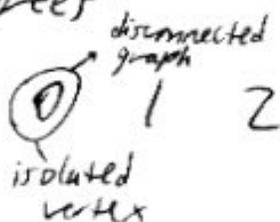
2) Induction

$$3) |E \text{ for } K_n| = \frac{n(n-1)}{2} \xrightarrow{\text{Replace } n \text{ w/ } n+1} \frac{(n+1)n}{2} = 1+2+\dots+n$$

"Old friend"

In a simple graph on $n (\geq 2)$ vertices, at least two vertices must have the same degree.

Otherwise, all n vertices must have different degrees



can't have both

Smells like pigeonhole principle

Old HW: 4.1.55

Convex:
Every line seg connecting two pts. in the interior or boundary of the polygon lies entirely within this set.

How many diagonals does a convex polygon with n sides have?

Ex



sides aren't diagonals

no



Each vertex has a diagonal to all vertices except itself and both neighbors.

\uparrow $n-3$ diagonals

$$\frac{n(n-3)}{2}$$

\leftarrow to adjust for double-counting

Weighted graphs (TSP)



Many applies

- Chemistry
- Scheduling (coloring vertices/edges)
- Circuit design
- Telecommunications/Internet
- Traffic Networks
- Garbage collection
- Social sciences

15-7 in math 21

4-color Thm: how many colors are needed to color a map

s.t. no 2 adjacent countries get the same color. (No Michigan)



& it's always enough! Countries - vertices
Adjacencies - edges

Described 1852 - Computer pt. 1976

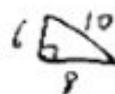
Political setup from Star Trek

Critical # Th'y Theorem:
Fermat's Last Theorem (Conjecture - 350 yrs. old)
Proven by Andrew Wiles of Princeton.

$$x^2 + y^2 = z^2$$

Find $x, y, z \in \mathbb{Z}^+$ that solve this.
There are many triples:

$$3^2 + 4^2 = 5^2$$



$$6^2 + 8^2 = 10^2$$

$$5^2 + 12^2 = 13^2$$

$$8^2 + 15^2 = 17^2$$

⋮

∞ many triples

Thm: for $n \in \mathbb{Z}, n \geq 3$

$$x^n + y^n = z^n$$

has no solution (x, y, z)
where $x, y, z \in \mathbb{Z}^+$

$$x^3 + y^3 = z^3$$

$$x^4 + y^4 = z^4$$

⋮

Fermat wrote in a margin

1983-
Gerd
Faltings
proved:
finitely
many
soln.

A Personal Favorite

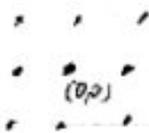
Consider the $\mathbb{Z} \times \mathbb{Z}$ lattice of pts.

set of "all" pts. (x, y)

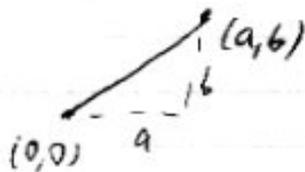
both integer



Is there a straight line that passes through $(0,0)$ and no other lattice point?



Let's assume we have a straight line that hits $(0,0)$ and some other lattice point (a,b) .



The slope must be $\frac{b}{a}$ $\exists \in \mathbb{Z}$ or undefined \downarrow
The slope must be rational ($\in \mathbb{Q}$)
or undefined.

We've shown

If a straight line hits $(0,0)$ and another lattice point (a,b) , then its slope is rational or undefined.

The converse must be true:

If the slope of a line is irrational, then the straight line can't hit both $(0,0)$ and another lattice pt.

Ex $y = \pi x$

Ex $y = \sqrt{2} x$

\vdots

Infinitely many lines, each corresp. to an irrat'l #.